

**600.363/463 Algorithms**  
**Mid Semester Examination 1**  
**October 11, 2013**  
**1 hr 10 mins; closed book**

I. Specify whether the following equations hold or not. Give an intuitive justification for your answer. Formal proofs are not needed.

1.  $3n^2 + 5n = O(n^2 \log n)$
2.  $2^n = O(n^4 \log \log n)$
3.  $10n^2 = O(n^2 - 100n)$
4.  $(n^2 + n)(n \log n + n^{1.1}) = O(n^3)$
5.  $n^3 \log n = \Theta(n^3)$
6.  $n^3 \log^5 n = O(n^{3.01})$
7.  $4^n = \Theta(3^n)$
8.  $n^2 \log n (\log \log n)^4 = O(n^2 \log^2 n)$
9.  $\log^2 n = O(n^{0.05})$
10.  $n^2 + n^{1.9} = O(n^2 + n^{0.5})$

II. Solve the following recurrences. You can invoke a known result and write down the answer, or you can justify the answer by repeated substitutions or by induction. When needed you can assume appropriate initial values.

1.  $T(n) \leq 3T(\frac{n}{2}) + n^2$
2.  $T(n) \leq 5T(\frac{n}{2}) + n^2$
3.  $T(n) \leq 2T(\frac{n}{4}) + T(\frac{n}{3})$
4.  $T(n) \leq 2nT(n-1)$
5.  $T(n) \leq 2T(n-1) + T(n-2)$

III. Let  $\delta$  be a positive-valued *affinity* function between pairs of symbols. Then for any pair of equal length strings  $X = a_1 a_2 \cdots a_\ell$  and  $Y = b_1 b_2 \cdots b_\ell$ ,  $\delta(X, Y) = \sum_{i=1}^{\ell} \delta(a_i, b_i)$ . For any 2 strings  $X = a_1 a_2 \cdots a_n$ ,  $Y = b_1 b_2 \cdots b_m$ ,  $n \leq m$ ,  $\delta(X, Y)$  is defined as

$$\max_{|Z|=n, Z \text{ a subsequence of } Y} \delta(X, Z)$$

For example, if  $\delta(a, a) = 2$ ,  $\delta(a, b) = \delta(b, a) = 1$ ,  $\delta(b, b) = 3$ , then  $\delta(abab, bbab) = 1+3+2+3 = 9$ ;  $\delta(abba, \underline{bbaaa}) = 7$ , and a subsequence that maximizes the value is underlined.

Design a dynamic programming based algorithm which for a given  $\delta$ ,  $X$  and  $Y$  computes  $\delta(X, Y)$ .

Hint: If  $X = x_1x_2 \cdots x_n$ ,  $Y = y_1y_2 \cdots y_m$ , for any  $i \leq j$ , let  $c[i, j] = \delta(x_1x_2 \cdots x_i, y_1y_2 \cdots y_j)$ .

- a) Specify a computation formula for  $c[i, j]$ ,
- b) Specify an order in which the entries of the  $c$  matrix can be computed, and
- c) What is the speed of the algorithm?

IV.

1. Assume that a set  $S$  of  $n$  numbers are stored in some form of balanced binary search tree; i.e. the depth of the tree is  $O(\log n)$ . In addition to the key value and the pointers to children, assume that every node contains the number of nodes in its subtree. Design  $O(\log n)$  step algorithms for performing the following operations.
  - (a) Given a positive integer  $k$ ,  $1 \leq k \leq n$ , compute the  $k^{\text{th}}$  smallest element of  $S$ , and
  - (b) Given 2 numbers  $x$  and  $y$ , compute the size of the subset  $\{z \mid z \in S, x \leq z \leq y\}$ .
2. Specify a reason why a balanced binary tree is better than a complete binary tree for storing the set  $S$ :
  - (a) A balanced binary tree is more space efficient,
  - (b) Inserts and Deletes can be performed faster, or
  - (c) Implementation of a balanced binary tree requires less RAM space