

## Efficient Parsing for

- Bilexical CF Grammars
- Head Automaton Grammars

Jason Eisner  
U. of Pennsylvania

Giorgio Satta  
U. of Padova, Italy



## When's a grammar bilexical?

If it has rules / entries that mention 2 specific words in a dependency relation:

convene - meeting  
eat - blintzes  
ball - bounces  
joust - with

## Bilexical Grammars

- Instead of  $VP \rightarrow V NP$
- or even  $VP \rightarrow \text{**solved**} NP$
- use detailed rules that mention **2 heads**:
  - $S[\text{**solved**}] \rightarrow NP[\text{**Peggy**}] VP[\text{**solved**}]$
  - $VP[\text{**solved**}] \rightarrow V[\text{**solved**}] NP[\text{**puzzle**}]$
  - $NP[\text{**puzzle**}] \rightarrow Det[\text{**a**}] N[\text{**puzzle**}]$
- so we can exclude, or reduce probability of,
  - $VP[\text{**solved**}] \rightarrow V[\text{**solved**}] NP[\text{**goat**}]$
  - $NP[\text{**puzzle**}] \rightarrow Det[\text{**two**}] N[\text{**puzzle**}]$

## Bilexical CF grammars

- Every rule has one of these forms:
  - $A[x] \rightarrow B[x] C[y]$  *so head of LHS*
  - $A[x] \rightarrow B[y] C[x]$  *is inherited from*
  - $A[x] \rightarrow x$  *a child on RHS.*(rules could also have probabilities)

$B[x], B[y], C[x], C[y], \dots$  many nonterminals  
 $A, B, C \dots$  are "traditional nonterminals"  
 $x, y \dots$  are words

## Bilexicalism at Work

- Not just selectional but adjunct preferences:
  - Peggy **solved** a puzzle **from** the library.
  - Peggy solved [a **puzzle from** the library].

Hindle & Rooth (1993) - PP attachment

## Bilexicalism at Work

Bilexical parsers that fit the CF formalism:

- Alshawi (1996) - head automata
- Charniak (1997) - Treebank grammars
- ★ Collins (1997) - context-free grammars
- Eisner (1996) - dependency grammars

Other superlexicalized parsers that don't:

- Jones & Eisner (1992) - bilexical LFG parser
- Lafferty et al. (1992) - stochastic link parsing
- Magerman (1995) - decision-tree parsing
- Ratnaparkhi (1997) - maximum entropy parsing
- Chelba & Jelinek (1998) - shift-reduce parsing

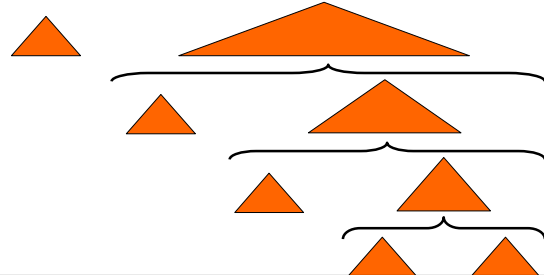
## How bad is bilex CF parsing?

$$A[x] \rightarrow B[x] C[y]$$

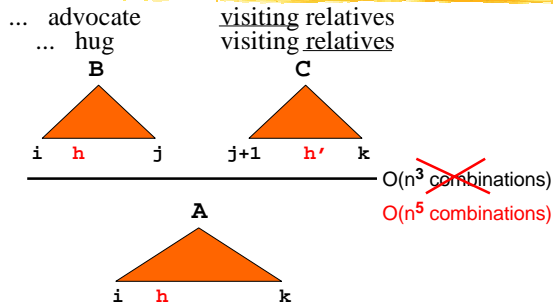
- Grammar size =  $O(t^3 V^2)$   
where  $t = |\{A, B, \dots\}|$   $V = |\{x, y, \dots\}|$
- So CKY takes  $O(t^3 V^2 n^3)$
- Reduce to  $O(t^3 n^5)$  since relevant  $V = n$
- This is terrible ... can we do better?
  - Recall: regular CKY is  $O(t^3 n^3)$

## The CKY-style algorithm

[Mary] → loves    [[the] → girl ← [outdoors]]

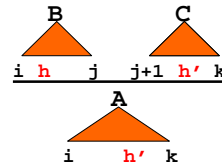


## Why CKY is $O(n^5)$ not $O(n^3)$

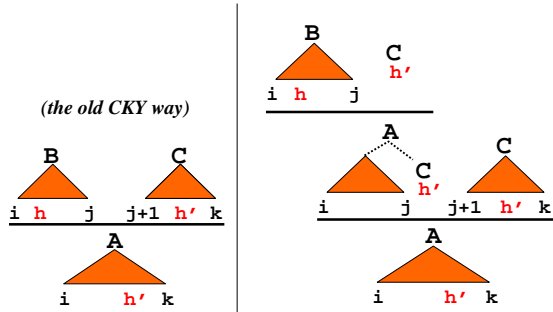


## Idea #1

- Combine B with what C?
- must try different-width C's (vary  $k$ )
- must try differently-headed C's (vary  $h'$ )
- Separate these!



## Idea #1



## Head Automaton Grammars

(Alshawi 1996)

[Good old **Peggy**] **solved** [the puzzle] [with her teeth] !

The **head automaton** for **solved**:

- a finite-state device
- can consume words adjacent to it on either side
- does so after they've consumed *their* dependents

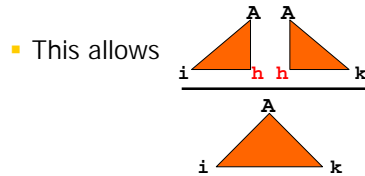
[Peggy] **solved** [puzzle] [with]    (state = V)  
 [Peggy] **solved** [with]                (state = VP)  
 [Peggy] **solved**                        (state = VP)  
    **solved**                                (state = S; halt)

## Formalisms too powerful?

- So we have Bilex CFG and HAG in  $O(n^4)$ .
- HAG is quite powerful - head  $c$  can require  $a^n \underline{c} b^n$ :  
 $\dots [ \dots a_3 \dots ] [ \dots a_2 \dots ] [ \dots a_1 \dots ] \underline{c} [ \dots b_1 \dots ] [ \dots b_2 \dots ] [ \dots b_3 \dots ] \dots$   
*not center-embedding*,  $[ a_3 [ [ a_2 [ [ a_1 ] b_1 ] ] b_2 ] ] b_3$ 
  - Linguistically unattested and unlikely
  - Possible only if the HA has a left-right cycle
  - Absent such cycles, can we parse faster?**
    - (for both HAG and equivalent Bilexical CFG)

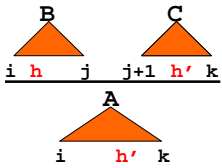
## Transform the grammar

- Absent such cycles, we can transform to a "split grammar":
  - Each head eats all its right dependents first
  - I.e., left dependents are more oblique.



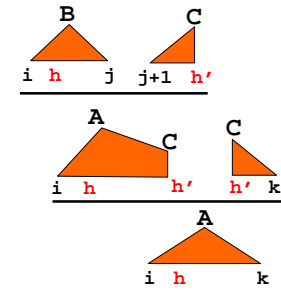
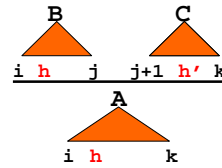
## Idea #2

- Combine what B and C?
  - must try different-width C's (vary  $k$ )
  - must try different midpoints  $j$
- Separate these!



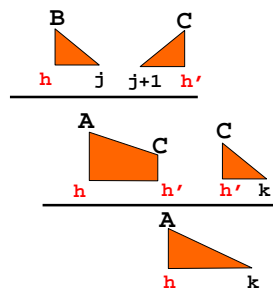
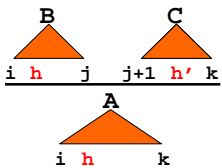
## Idea #2

(the old CKY way)



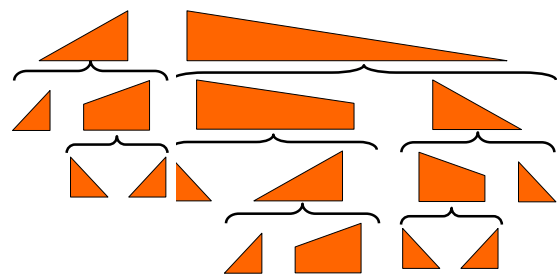
## Idea #2

(the old CKY way)



## The $O(n^3)$ half-tree algorithm

[Mary] → loves [[the] → girl ← [outdoors]]



## Theoretical Speedup

- $n$  = input length       $g$  = polysemy
- $t$  = traditional nonterms or automaton states
- Naive:  $O(n^5 g^2 t)$
- **New:  $O(n^4 g^2 t)$**
- Even better for split grammars:
  - Eisner (1997):  $O(n^3 g^3 t^2)$
  - **New:  $O(n^3 g^2 t)$**   
*all independent of vocabulary size!*

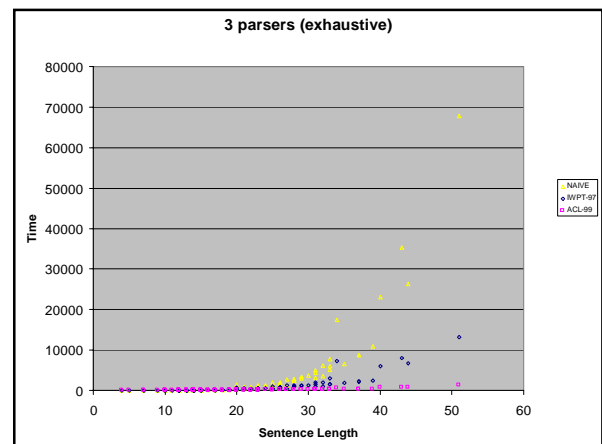
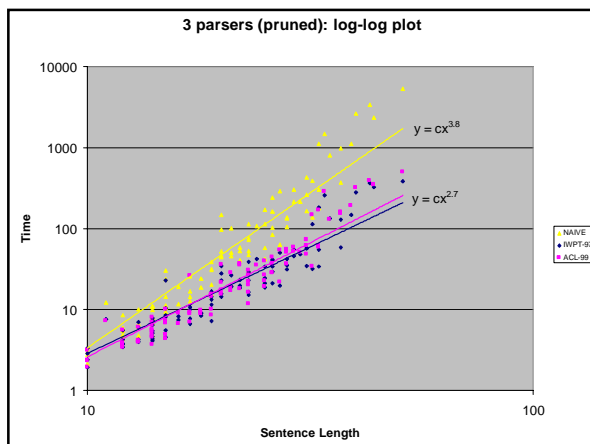
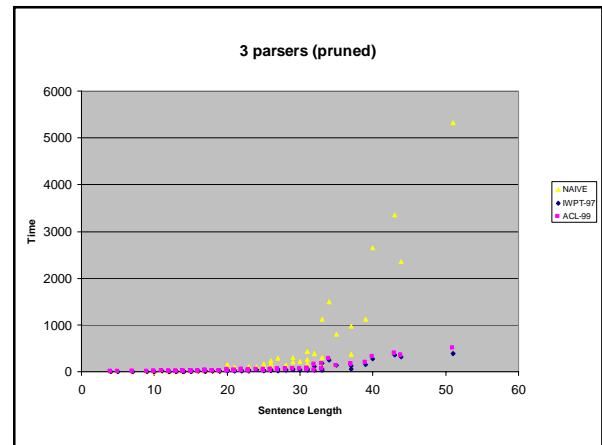
## Reality check

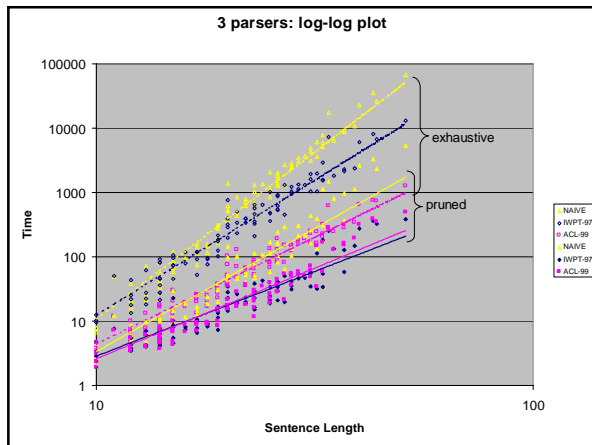
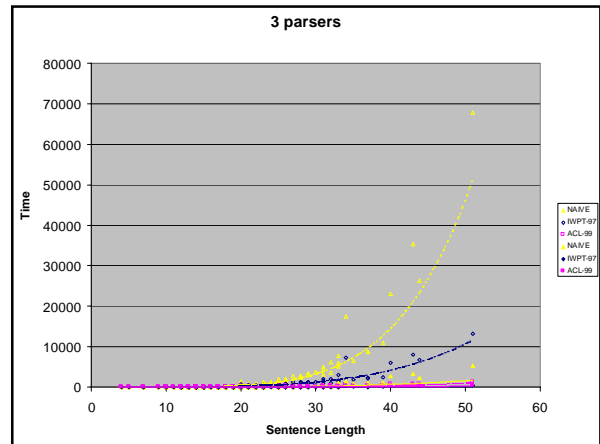
- Constant factor
- Pruning may do just as well
  - "visiting relatives": 2 plausible NP hypotheses
  - i.e., both heads survive to compete - common??
- Amdahl's law
  - much of time spent smoothing probabilities
  - fixed cost per parse if we cache probs for reuse

## Experimental Speedup (not in paper)

Used Eisner (1996) Treebank WSJ parser and its split bilinear grammar

- Parsing with pruning:
  - Both old and new  $O(n^3)$  methods give **5x** speedup over the  $O(n^5)$  - at 30 words
- Exhaustive parsing (e.g., for EM):
  - Old  $O(n^3)$  method (Eisner 1997) gave **3x** speedup over  $O(n^5)$  - at 30 words
  - **New  $O(n^3)$  method gives 19x speedup**





## Summary

- Simple bilexical CFG notion  $A[x] \rightarrow B[x] C[y]$
- Covers several existing stat NLP parsers
- Fully general  $O(n^4)$  algorithm - not  $O(n^5)$
- Faster  $O(n^3)$  algorithm for the "split" case
- Demonstrated practical speedup
- Extensions*: TAGs and post-transductions